

TMS416160, TMS416160P 1 048 576-WORD BY 16-BIT HIGH-SPEED DYNAMIC RANDOM-ACCESS MEMORIES

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- Organization . . . 1 048 576 x 16
- Single 5-V Supply (10% Tolerance)
- Performance Ranges:

	ACCESS TIME t _{TRAC} MAX	ACCESS TIME t _{CAC} MAX	ACCESS TIME t _{AA} MAX	READ OR WRITE CYCLE MIN
'416160/P-60	60 ns	15 ns	30 ns	110 ns
'416160/P-70	70 ns	18 ns	35 ns	130 ns
'416160/P-80	80 ns	20 ns	40 ns	150 ns

- Enhanced Page Mode Operation With CAS-Before-RAS Refresh
- Long Refresh Period . . .
4096-Cycle Refresh in 64 ms (Max)
512 ms Max for Low-Power, Self-Refresh Version (TMS416160P)
- 3-State Unlatched Output
- Low Power Dissipation
- Self-Refresh with Low Power
- All Inputs, Outputs, and Clocks are TTL Compatible
- High-Reliability Plastic 42-Lead 400-Mil-Wide Surface Mount (SOJ) Package, and 44/50-Lead Thin Small Outline Package (TSOP)
- Operating Free-Air Temperature Range 0°C to 70°C
- Texas Instruments EPIC™ CMOS Process

description

The TMS416160 series are high-speed, 16 777 216-bit dynamic random-access memories organized as 1 048 576 words of sixteen bits each.

The TMS416160P series are high-speed, low-power, self-refresh, 16 777 216-bit dynamic random-access memories organized as 1 048 576 words of sixteen bits each.

They employ state-of-the-art EPIC™ (Enhanced Performance Implanted CMOS) technology for high performance, reliability, and low power at low cost.

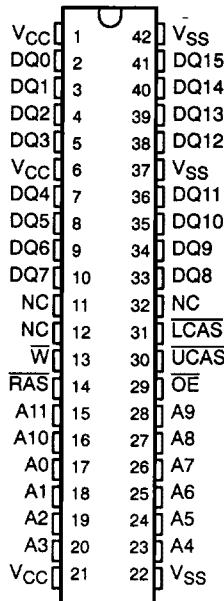
These devices feature maximum $\overline{\text{RAS}}$ access times of 60 ns, 70 ns, and 80 ns. Maximum power dissipation is as low as 385 mW operating and 11 mW standby on 80 ns devices.

All inputs and outputs, including clocks, are compatible with Series 74 TTL. All addresses and data-in lines are latched on-chip to simplify system design. Data out is unlatched to allow greater system flexibility.

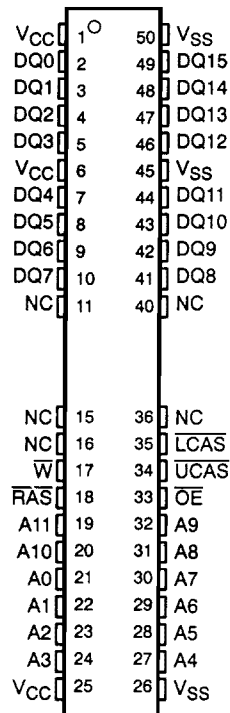
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RE PACKAGE†
(TOP VIEW)



DC PACKAGE†
(TOP VIEW)



† Packages are shown for pinout reference only.

PIN NOMENCLATURE	
A0-A11	Address Inputs
DQ0-DQ15	Data In/Data Out
$\overline{\text{LCAS}}$	Lower Column-Address Strobe
$\overline{\text{UCAS}}$	Upper Column-Address Strobe
NC	No Internal Connection
$\overline{\text{OE}}$	Output Enable
$\overline{\text{RAS}}$	Row-Address Strobe
$\overline{\text{W}}$	Write Enable
VCC	5-V Supply
VSS	Ground

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The TMS416160 and TMS416160P are each offered in a 42-lead plastic surface mount SOJ (RE suffix) package, and a 44/50-lead plastic surface mount TSOP (DC suffix). These packages are characterized for operation from 0°C to 70°C.

operation

dual $\overline{\text{CAS}}$

Two $\overline{\text{CAS}}$ pins ($\overline{\text{LCAS}}$ – $\overline{\text{UCAS}}$) are provided to give independent control of the sixteen data I/O pins (DQ0–DQ15), with $\overline{\text{LCAS}}$ corresponding to DQ0–DQ7 and $\overline{\text{UCAS}}$ corresponding to DQ8–DQ15. For read or write cycles, the column address is latched on the first $\overline{\text{xCAS}}$ falling edge. Each $\overline{\text{xCAS}}$ pin going low enables its corresponding DQ pin with data coming from the column address to be latched on the first falling $\overline{\text{xCAS}}$ edge. All address setup and hold parameters are referenced to the first falling $\overline{\text{xCAS}}$ edge. The delay time from $\overline{\text{xCAS}}$ low to valid data out (see parameter t_{CAC}) is measured from each individual $\overline{\text{CAS}}$ to its corresponding DQx pin.

In order to latch in a new column address, all $\overline{\text{xCAS}}$ pins must be brought high. The column precharge time (see parameter t_{CP}) is measured from the last $\overline{\text{xCAS}}$ rising edge to the first falling $\overline{\text{xCAS}}$ edge of the new cycle. Keeping a column address valid while toggling $\overline{\text{xCAS}}$ requires a minimum setup time, t_{CLCH} . During t_{CLCH} , at least one $\overline{\text{xCAS}}$ must be brought low before the other $\overline{\text{xCAS}}$ is taken high.

For early write cycles, the data is latched on the first falling $\overline{\text{xCAS}}$ edge. Only the DQs that have the corresponding $\overline{\text{xCAS}}$ low will be written into. Each $\overline{\text{xCAS}}$ will have to meet t_{CAS} minimum in order to ensure writing into the storage cell. In order to latch a new address and new data, all $\overline{\text{xCAS}}$ pins need to come high and meet t_{CP} .

enhanced page mode

Page-mode operation allows faster memory access by keeping the same row address while selecting random column addresses. The time for row-address setup and hold and address multiplex is thus eliminated. The maximum number of columns that may be accessed is determined by the maximum $\overline{\text{RAS}}$ low time and the $\overline{\text{xCAS}}$ page-mode cycle time used. With minimum $\overline{\text{xCAS}}$ page cycle time, all 256 columns specified by column addresses A0 through A7 can be accessed without intervening $\overline{\text{RAS}}$ cycles.

Unlike conventional page-mode DRAMs, the column-address buffers in this device are activated on the falling edge of $\overline{\text{RAS}}$. The buffers act as transparent or flow-through latches while $\overline{\text{xCAS}}$ is high. The falling edge of the first $\overline{\text{xCAS}}$ latches the column addresses. This feature allows the devices to operate at a higher data bandwidth than conventional page-mode parts, because data retrieval begins as soon as column address is valid rather than when $\overline{\text{xCAS}}$ transitions low. This performance improvement is referred to as *enhanced page mode*. Valid column address may be presented immediately after t_{RAH} (row address hold time) has been satisfied, usually well in advance of the falling edge of $\overline{\text{xCAS}}$. In this case, data is obtained after $t_{\text{CAC max}}$ (access time from $\overline{\text{xCAS}}$ low) if $t_{\text{AA max}}$ (access time from column address) has been satisfied. In the event that column addresses for the next page cycle are valid at the time $\overline{\text{xCAS}}$ goes high, minimum access time for the next cycle is determined by t_{CPA} (access time from rising edge of the last $\overline{\text{xCAS}}$).

address (A0–A11)

Twenty address bits are required to decode 1 of 1 048 576 storage cell locations. Twelve row-address bits are set up on pins A0 through A11 and latched onto the chip by $\overline{\text{RAS}}$. Then, eight column-address bits are set up on pins A0 through A7 and latched onto the chip by the first $\overline{\text{xCAS}}$. All addresses must be stable on or before the falling edge of $\overline{\text{RAS}}$ and $\overline{\text{xCAS}}$. $\overline{\text{RAS}}$ is similar to a chip enable in that it activates the sense amplifiers as well as the row decoder. $\overline{\text{xCAS}}$ is used as a chip select, activating its corresponding output buffer and latching the address bits into the column-address buffers.

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write enable (\overline{W})

The read or write mode is selected through the \overline{W} input. A logic high on the \overline{W} input selects the read mode and a logic low selects the write mode. The write-enable terminal can be driven from the standard TTL circuits without a pullup resistor. The data inputs are disabled when the read mode is selected. When \overline{W} goes low prior to \overline{xCAS} (early write), data out will remain in the high-impedance state for the entire cycle permitting a write operation with \overline{OE} grounded.

data in (DQ0–DQ15)

Data is written during a write or read-modify-write cycle. Depending on the mode of operation, the falling edge of \overline{xCAS} or \overline{W} strobes data into the on-chip data latch. In an early-write cycle, \overline{W} is brought low prior to \overline{xCAS} and the data is strobed in by the first occurring \overline{xCAS} with setup and hold times referenced to this signal. In a delayed-write or read-modify-write cycle, \overline{xCAS} will already be low, thus the data will be strobed in by \overline{W} with setup and hold times referenced to this signal. In a delayed-write or read-modify-write cycle, \overline{OE} must be high to bring the output buffers to high impedance prior to impressing data on the I/O lines.

data out (DQ0–DQ15)

The three-state output buffer provides direct TTL compatibility (no pullup resistor required) with a fanout of two Series 74 TTL loads. Data out is the same polarity as data in. The output is in the high-impedance (floating) state until \overline{xCAS} and \overline{OE} are brought low. In a read cycle, the output becomes valid after the access time interval t_{CAS} (which begins with the negative transition of \overline{xCAS}) as long as t_{RAC} and t_{AA} are satisfied.

output enable (\overline{OE})

\overline{OE} controls the impedance of the output buffers. When \overline{OE} is high, the buffers will remain in the high-impedance state. Bringing \overline{OE} low during a normal cycle will activate the output buffers, putting them in the low-impedance state. It is necessary for both \overline{RAS} and \overline{xCAS} to be brought low for the output buffers to go into low-impedance state, they will remain in the low-impedance state until either \overline{OE} or \overline{xCAS} is brought high.

\overline{RAS} -only refresh

A refresh operation must be performed at least once every sixty-four milliseconds (512 ms for TMS416160P) to retain data. This can be achieved by strobing each of the 4096 rows (A0–A11). A normal read or write cycle will refresh all bits in each row that is selected. A \overline{RAS} -only operation can be used by holding all \overline{xCAS} at the high (inactive) level, thus conserving power as the output buffers remain in the high-impedance state. Externally generated addresses must be used for a \overline{RAS} -only refresh.

hidden refresh

Hidden refresh may be performed while maintaining valid data at the output pin. This is accomplished by holding \overline{xCAS} at V_{IL} after a read operation and cycling \overline{RAS} after a specified precharge period, similar to a \overline{RAS} -only refresh cycle.

\overline{xCAS} -before- \overline{RAS} refresh

\overline{xCAS} -before- \overline{RAS} refresh is utilized by bringing at least one \overline{xCAS} low earlier than \overline{RAS} (see parameter t_{CSR}) and holding it low after \overline{RAS} falls (see parameter t_{CHR}). For successive \overline{xCAS} -before- \overline{RAS} refresh cycles, \overline{xCAS} can remain low while cycling \overline{RAS} . The external address is ignored and the refresh address is generated internally. The external address is also ignored during the hidden refresh option.

A low-power battery-backup refresh mode that requires less than 500 μA refresh current is available on the TMS416160P. Data integrity is maintained using \overline{xCAS} -before- \overline{RAS} refresh with a period of 125 μs while holding \overline{RAS} low for less than 1 μs . To minimize current consumption, all input levels need to be at CMOS levels ($V_{IL} < 0.2 V$, $V_{IH} > V_{CC} - 0.2 V$).

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self refresh (TMS416160P)

The self-refresh mode is entered by dropping \overline{xCAS} low prior to \overline{RAS} going low. Then \overline{xCAS} and \overline{RAS} are both held low for a minimum of 100 μs . The chip is then refreshed internally by an on-board oscillator. No external address is required since the CBR counter is used to keep track of the address. To exit the self-refresh mode both \overline{RAS} and \overline{xCAS} are brought high to satisfy t_{CHS} . Upon exiting self-refresh mode, a burst refresh (refresh a full set of row addresses) must be executed before continuing with normal operation. The burst refresh ensures the DRAM is fully refreshed.

power up

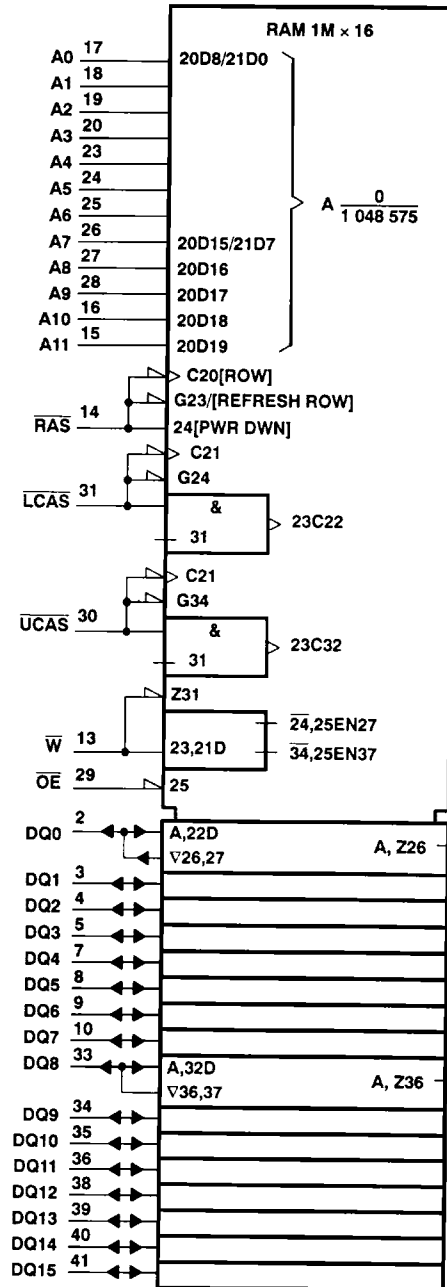
To achieve proper device operation, an initial pause of 200 μs followed by a minimum of eight \overline{RAS} cycles is required after power up to the full V_{CC} level.

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logic symbol†

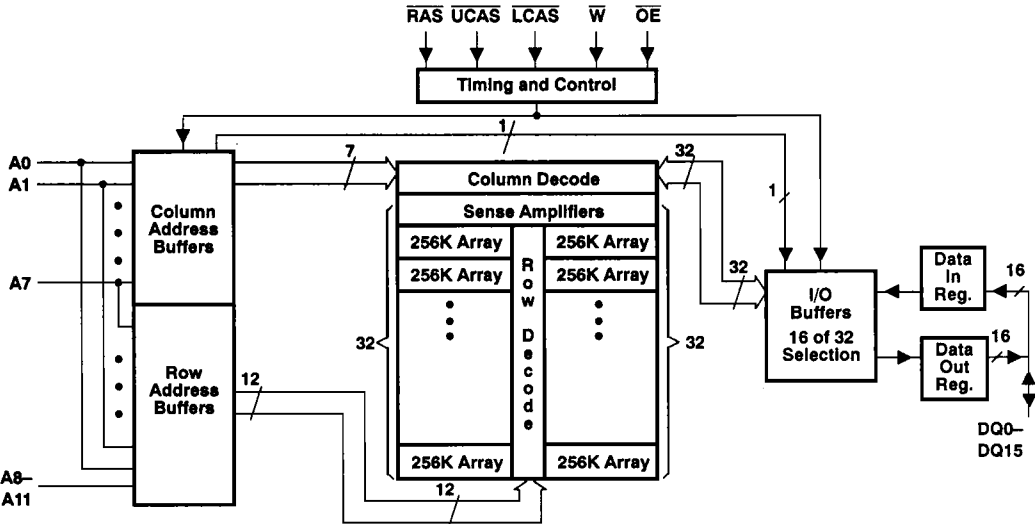


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† This symbol is in accordance with ANSI/IEEE Std 91-1984 and IEC Publication 617-12.
 The pin numbers shown correspond to the RE package.



functional block diagram



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absolute maximum ratings over operating free-air temperature range (unless otherwise noted)†

- Supply voltage range on any pin (see Note 1) - 1 V to 7 V
- Supply voltage range on V_{CC} - 1 V to 7 V
- Short circuit output current 50 mA
- Power dissipation 1 W
- Operating free-air temperature range 0°C to 70°C
- Storage temperature range - 55°C to 125°C

† Stresses beyond those listed under "absolute maximum ratings" may cause permanent damage to the device. This is a stress rating only, and functional operation of the device at these or any other conditions beyond those indicated in the "recommended operating conditions" section of this specification is not implied. Exposure to absolute-maximum-rated conditions for extended periods may affect device reliability.

NOTE 1: All voltage values in this data sheet are with respect to V_{SS}.

recommended operating conditions

	MIN	NOM	MAX	UNIT
V _{CC} Supply voltage	4.5	5	5.5	V
V _{SS} Supply voltage		0		V
V _{IH} High-level input voltage	2.4		6.5	V
V _{IL} Low-level input voltage (see Note 2)	-1		0.8	V
T _A Operating free-air temperature	0		70	°C

NOTE 2: The algebraic convention, where the more negative (less positive) limit is designated as minimum, is used in this data sheet for logic voltage levels only.

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electrical characteristics over full ranges of recommended operating conditions (unless otherwise noted)

PARAMETER	TEST CONDITIONS	'416160-60 '416160P-60		'416160-70 '416160P-70		'416160-80 '416160P-80		UNIT
		MIN	MAX	MIN	MAX	MIN	MAX	
V _{OH}	High-level output voltage I _{OH} = - 5 mA	2.4		2.4		2.4		V
V _{OL}	Low-level output voltage I _{OL} = 4.2 mA	0.4		0.4		0.4		V
I _I	Input current (leakage) V _{CC} = 5.5 V, V _I = 0 to 6.5 V, All other pins = 0 V to V _{CC}	± 10		± 10		± 10		µA
I _O	Output current (leakage) V _{CC} = 5.5 V, V _O = 0 to V _{CC} , $\overline{\text{xCAS}}$ high	± 10		± 10		± 10		µA
I _{CC1} †‡	Read or write cycle current V _{CC} = 5.5 V, Minimum cycle	90		80		70		mA
I _{CC2}	Standby current V _{IH} = 2.4 V (TTL), After 1 memory cycle, $\overline{\text{RAS}}$ and $\overline{\text{xCAS}}$ high			2		2		mA
		V _{IH} = V _{CC} - 0.2 V (CMOS), After 1 memory cycle, $\overline{\text{RAS}}$ and $\overline{\text{xCAS}}$ high	'416160	1	1	1	mA	
		'416160P	500	500	500	µA		
I _{CC3} ‡	Average refresh current ($\overline{\text{RAS}}$ -only or CBR) V _{CC} = 5.5 V, Minimum cycle, $\overline{\text{RAS}}$ cycling, $\overline{\text{xCAS}}$ high ($\overline{\text{RAS}}$ only) $\overline{\text{RAS}}$ low after $\overline{\text{xCAS}}$ low (CBR)	90		80		70		mA
I _{CC4} †§	Average page current V _{CC} = 5.5 V, t _{PC} = minimum, $\overline{\text{RAS}}$ low, $\overline{\text{xCAS}}$ cycling	90		80		70		mA
I _{CC6} ¶	Self refresh $\overline{\text{CAS}} < 0.2$ V, $\overline{\text{RAS}} < 0.2$ V, Measured after t _{RASS} minimum	500		500		500		µA
I _{CC7} †	Standby current, outputs enabled $\overline{\text{RAS}} = V_{IH}$, $\overline{\text{xCAS}} = V_{IL}$, Data out = enabled	5		5		5		mA
I _{CC10} ¶	Battery back-up operating current (equivalent refresh time is 512 ms). CBR only. t _{RC} = 125 µs, t _{RAS} ≤ 1 µs, V _{CC} - 0.2 V ≤ V _{IH} ≤ 6.5 V, 0 V ≤ V _{IL} ≤ 0.2 V, $\overline{\text{W}}$ and $\overline{\text{OE}} = V_{IH}$, Address and Data stable	500		500		500		µA

† Measured with outputs open.

‡ Measured with a maximum of one address change while $\overline{\text{RAS}} = V_{IL}$.

§ Measured with a maximum of one address change while $\overline{\text{xCAS}} = V_{IH}$.

¶ For TMS416160P only.

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capacitance over recommended ranges of supply voltage and operating free-air temperature,
f = 1 MHz (see Note 3)

PARAMETER		MIN	TYP	MAX	UNIT
C _{i(A)}	Input capacitance, address inputs			5	pF
C _{i(OE)}	Input capacitance, output enable			7	pF
C _{i(RC)}	Input capacitance, strobe inputs			7	pF
C _{i(W)}	Input capacitance, write-enable input			7	pF
C _O	Output capacitance			7	pF

NOTE 3: V_{CC} equal to 5 V ± 0.5 V and the bias on pins under test is 0 V.

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switching characteristics over recommended ranges of supply voltage and operating free-air temperature

PARAMETER	'416160-60 '416160P-60		'416160-70 '416160P-70		'416160-80 '416160P-80		UNIT
	MIN	MAX	MIN	MAX	MIN	MAX	
t _{CAC} Access time from xCAS low	15		18		20		ns
t _{AA} Access time from column address	30		35		40		ns
t _{RAC} Access time from RAS low	60		70		80		ns
t _{OEa} Access time from OE low	15		18		20		ns
t _{CPA} Access time from column precharge	35		40		45		ns
t _{CLZ} Delay time, xCAS low to output in low Z	0		0		0		ns
t _{OH} Output data hold time (from xCAS)	3		3		3		ns
t _{OHO} Output data hold time (from OE)	3		3		3		ns
t _{OFF} Output disable time after xCAS high (see Note 4)	0	15	0	18	0	20	ns
t _{OEZ} Output disable time after OE high (see Note 4)	0	15	0	18	0	20	ns

NOTE 4: t_{OFF} and t_{OEZ} are specified when the output is no longer driven.

timing requirements over recommended ranges of supply voltage and operating free-air temperature (see Note 5)

	'416160-60 '416160P-60		'416160-70 '416160P-70		'416160-80 '416160P-80		UNIT
	MIN	MAX	MIN	MAX	MIN	MAX	
t _{RC} Read cycle time (see Note 6)	110		130		150		ns
t _{WC} Write cycle time	110		130		150		ns
t _{RWC} Read-write/read-modify-write cycle time	155		181		205		ns
t _{PRWC} Page-mode read-modify-write cycle time	85		96		105		ns
t _{RASP} Page-mode pulse duration, RAS low (see Note 8)	60	100 000	70	100 000	80	100 000	ns
t _{RAS} Non-page-mode pulse duration, RAS low (see Note 8)	60	10 000	70	10 000	80	10 000	ns
t _{CAS} Pulse duration, xCAS low (see Note 9)	15	10 000	18	10 000	20	10 000	ns
t _{CP} Pulse duration, xCAS high (precharge)	10		10		10		ns
t _{RP} Pulse duration, RAS high (precharge)	40		50		60		ns
t _{WP} Write pulse duration	15		15		15		ns
t _{ASC} Column-address setup time before xCAS low	0		0		0		ns
t _{ASR} Row-address setup time before RAS low	0		0		0		ns
t _{DS} Data setup time before W low (see Note 10)	0		0		0		ns
t _{RCS} Read setup time before xCAS low	0		0		0		ns
t _{CWL} W-low setup time before xCAS high	15		18		20		ns

NOTES: 5. Timing measurements are referenced to V_{IL} max and V_{IH} min.

6. All cycle times assume t_T = 5 ns.

7. t_{PC} > t_{CP} min + t_{CAS} min + 2t_T.

8. In a read-modify-write cycle, t_{RWD} and t_{RWL} must be observed. Depending on the user's transition times, this may require additional RAS low time (t_{RAS}).

9. In a read-modify-write cycle, t_{CWD} and t_{CWL} must be observed. Depending on the user's transition times, this may require additional xCAS low time (t_{CAS}).

10. Reference to the first xCAS or W, whichever occurs last.

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timing requirements over recommended ranges of supply voltage and operating free-air temperature (continued) (see Note 5)

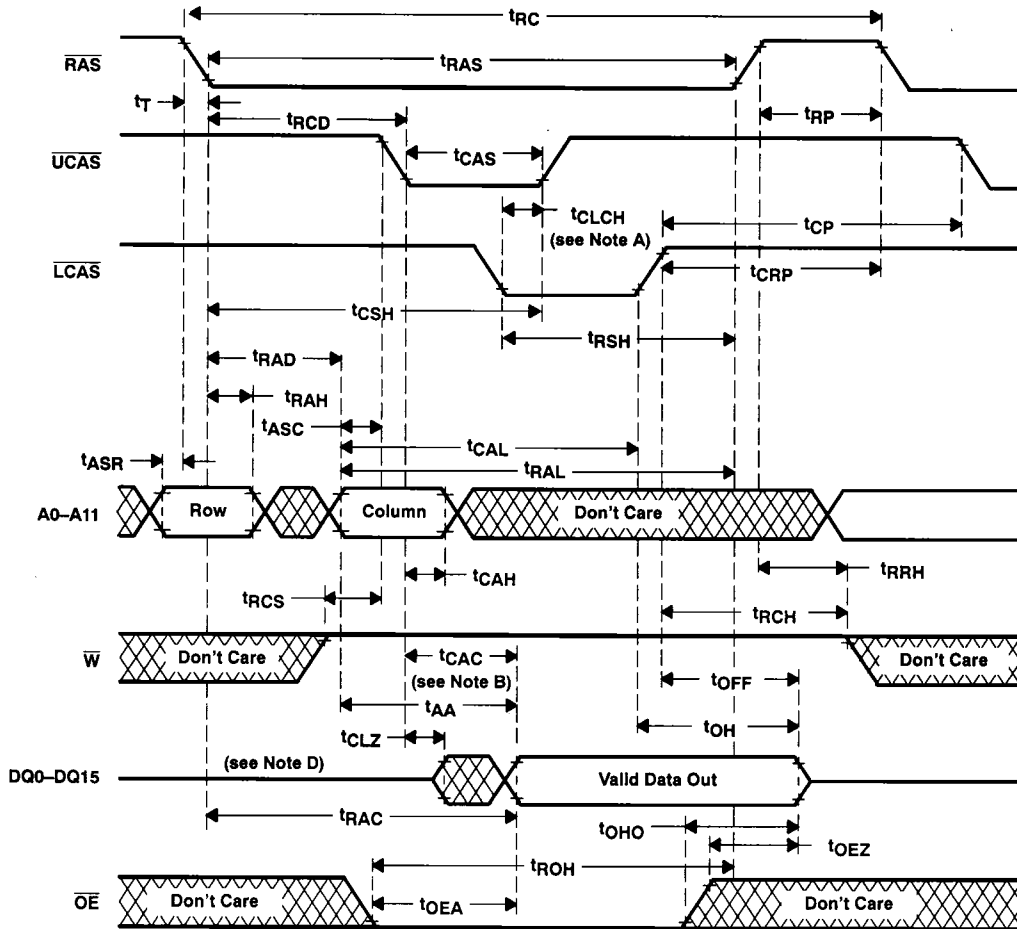
	'416160-60 '416160P-60		'416160-70 '416160P-70		'416160-80 '416160P-80		UNIT
	MIN	MAX	MIN	MAX	MIN	MAX	
t _{RWL} \bar{W} -low setup time before \bar{RAS} high	15		18		20		ns
t _{WCS} \bar{W} -low setup time before \bar{xCAS} low	0		0		0		ns
t _{CAH} Column-address hold time after \bar{xCAS} low	10		15		15		ns
t _{DH} Data hold time after \bar{xCAS} low (see Note 10)	10		15		15		ns
t _{RAH} Row-address hold time after \bar{RAS} low	10		10		10		ns
t _{RCH} Read hold time after \bar{xCAS} high (see Note 13)	0		0		0		ns
t _{RRH} Read hold time after \bar{RAS} high (see Note 13)	5		5		5		ns
t _{WCH} Write hold time after \bar{xCAS} low (see Note 12)	15		15		15		ns
t _{CLCH} Hold time, \bar{xCAS} low to \bar{xCAS} high	5		5		5		ns
t _{AWD} Delay time, column address to \bar{W} low (see Note 14)	55		63		70		ns
t _{CHR} Delay time, \bar{RAS} low to \bar{xCAS} high (see Note 11)	20		20		20		ns
t _{CRP} Delay time, \bar{xCAS} high to \bar{RAS} low	5		5		5		ns
t _{CSH} Delay time, \bar{RAS} low to \bar{xCAS} high	60		70		80		ns
t _{CSR} Delay time, \bar{xCAS} low to \bar{RAS} low (see Note 11)	10		10		10		ns
t _{CWD} Delay time, \bar{xCAS} low to \bar{W} low (see Note 14)	40		46		50		ns
t _{OEH} \bar{OE} command hold time	15		18		20		ns
t _{OED} Delay time, \bar{OE} high before data at DQ	15		18		20		ns
t _{ROH} Delay time, \bar{OE} low to \bar{RAS} high	10		10		10		ns
t _{RAD} Delay time, \bar{RAS} low to column address (see Note 15)	15	30	15	35	15	40	ns
t _{RAL} Delay time, column address to \bar{RAS} high	30		35		40		ns
t _{CAL} Delay time, column address to \bar{xCAS} high	30		35		40		ns
t _{RCD} Delay time, \bar{RAS} low to \bar{xCAS} low (see Note 15)	20	45	20	52	20	60	ns
t _{RPC} Delay time, \bar{RAS} high to \bar{xCAS} low	0		0		0		ns
t _{RSH} Delay time, \bar{xCAS} low to \bar{RAS} high	15		18		20		ns
t _{RWD} Delay time, \bar{RAS} low to \bar{W} low (see Note 14)	85		98		110		ns
t _{CPW} Delay time, \bar{W} from \bar{xCAS} precharge	60		68		75		ns
t _{CPRH} \bar{RAS} hold time from \bar{xCAS} precharge	35		40		45		ns
t _{CPR} \bar{xCAS} precharge before self refresh	0		0		0		ns
t _{RPS} \bar{RAS} precharge after self refresh	110		130		150		ns

- NOTES: 5. Timing measurements are referenced to V_{IL} max and V_{IH} min.
10. Reference to the first \bar{xCAS} or \bar{W} , whichever occurs last.
11. \bar{xCAS} -before- \bar{RAS} refresh only.
12. Early write operation only.
13. Either t_{RRH} or t_{RCH} must be satisfied for a read cycle.
14. Read-modify-write operation only.
15. Maximum value specified only to assure access time.

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PARAMETER MEASUREMENT INFORMATION

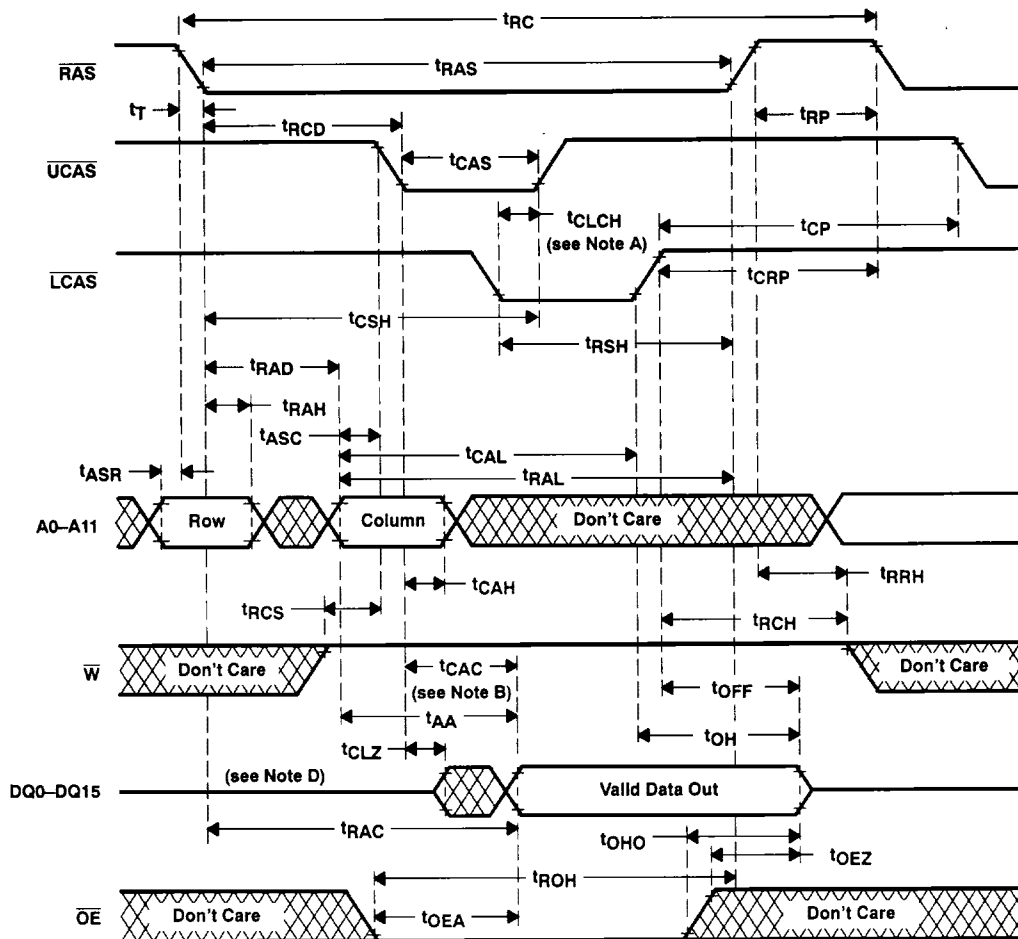


- NOTES: A. In order to hold the address latched by the first $\overline{x}CAS$ going low, the parameter t_{CLCH} must be met.
 B. t_{CAC} is measured from $\overline{x}CAS$ to its corresponding DQx.
 C. $\overline{x}CAS$ order is arbitrary.
 D. Output may go from high-impedance to an invalid data state prior to the specified access time.

Figure 2. Read Cycle Timing

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PARAMETER MEASUREMENT INFORMATION

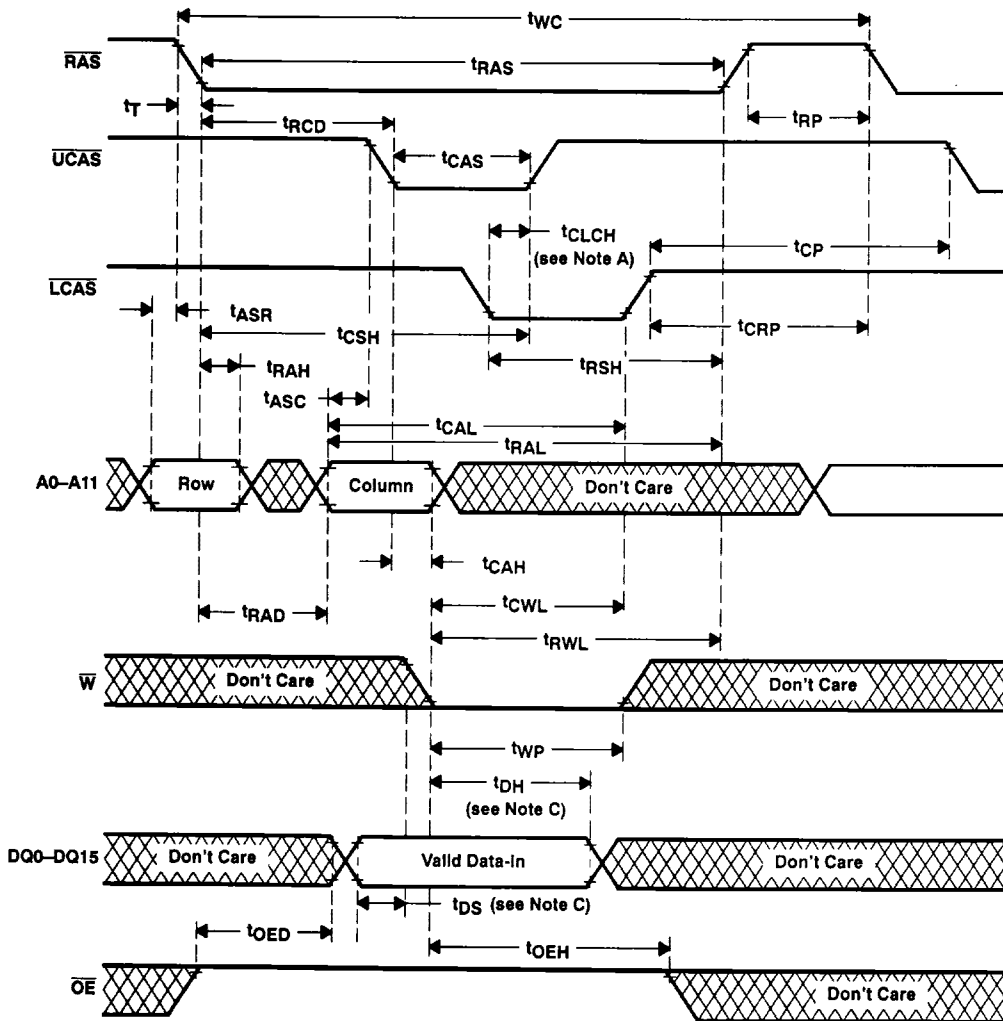


- NOTES: A. In order to hold the address latched by the first $\overline{x}CAS$ going low, the parameter t_{CLCH} must be met.
 B. t_{CAC} is measured from $\overline{x}CAS$ to its corresponding DQx.
 C. $\overline{x}CAS$ order is arbitrary.
 D. Output may go from high-impedance to an invalid data state prior to the specified access time.

Figure 2. Read Cycle Timing

PRODUCT PREVIEW

PARAMETER MEASUREMENT INFORMATION

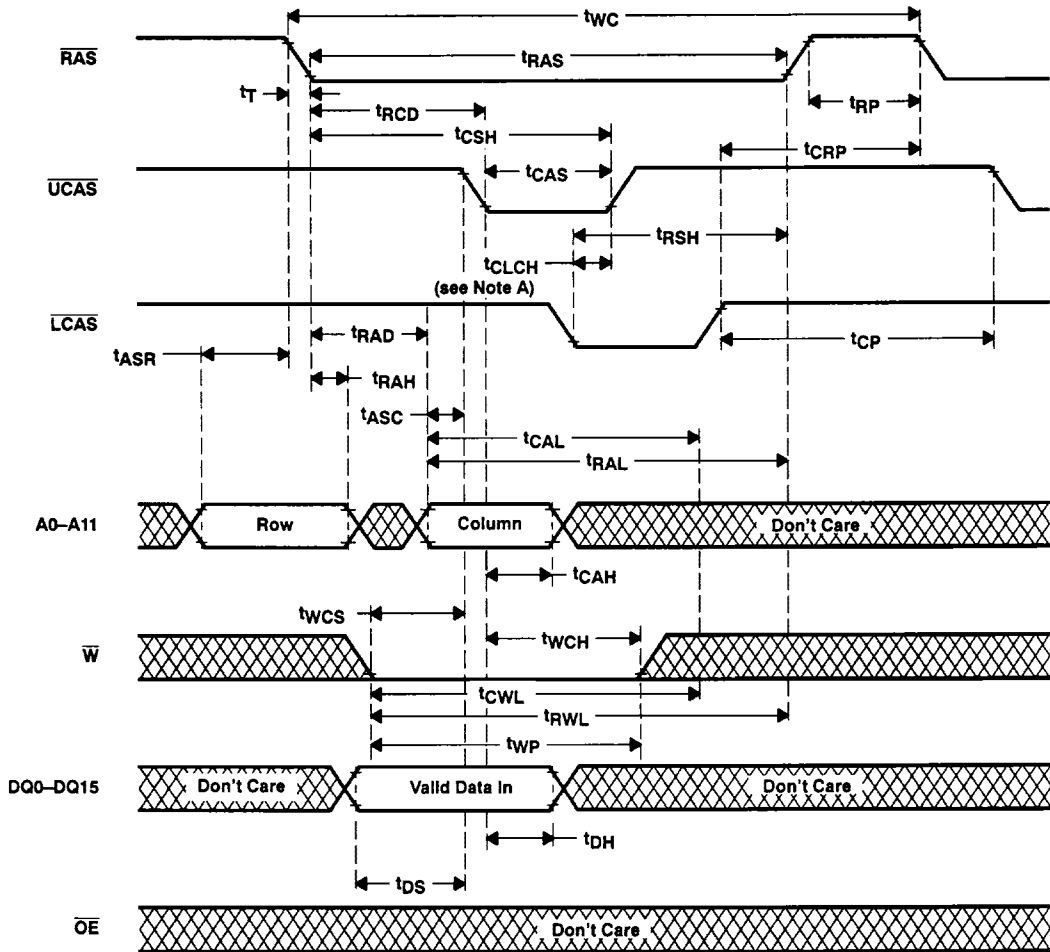


- NOTES: A. In order to hold the address latched by the first $\overline{x}CAS$ going low, the parameter t_{CLCH} must be met.
 B. $\overline{x}CAS$ order is arbitrary.
 C. Reference to the first $\overline{x}CAS$ or \overline{W} , whichever occurs last.

Figure 3. Write Cycle Timing

PRODUCT PREVIEW

PARAMETER MEASUREMENT INFORMATION

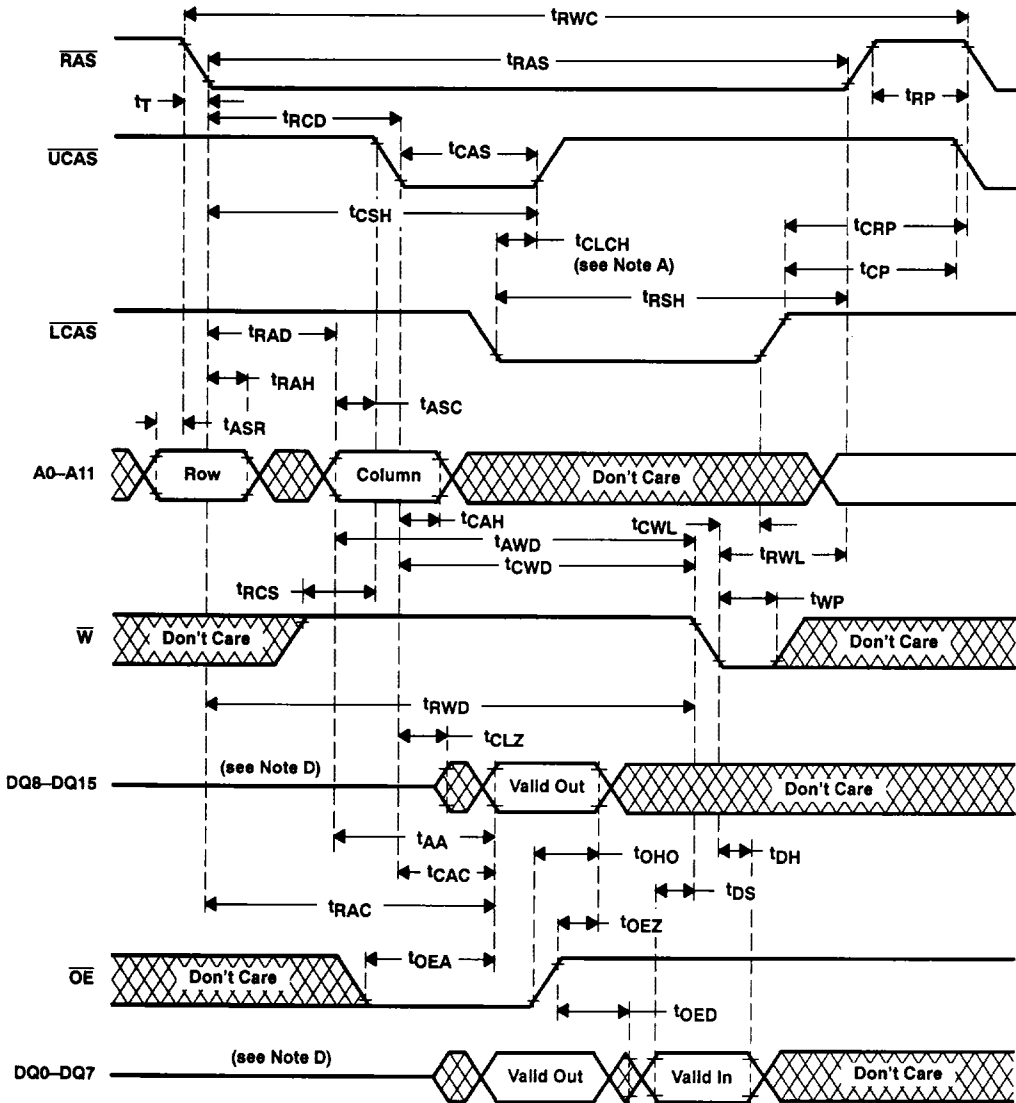


PRODUCT PREVIEW

- NOTES: A. In order to hold the address latched by the first $\overline{x}CAS$ going low, the parameter t_{CLCH} must be met.
 B. $\overline{x}CAS$ order is arbitrary.

Figure 4. Early Write Cycle Timing

PARAMETER MEASUREMENT INFORMATION

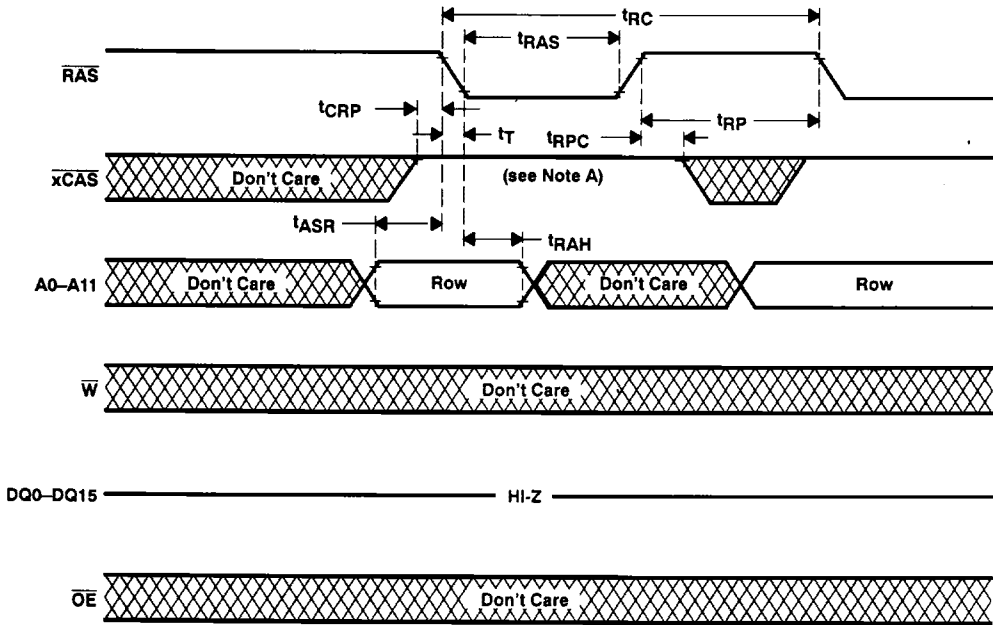


- NOTES: A. In order to hold the address latched by the first \overline{xCAS} going low, the parameter t_{CLCH} must be met.
 B. \overline{xCAS} order is arbitrary.
 C. t_{CAC} is measured from \overline{xCAS} to its corresponding DQx.
 D. Output might go from a high-impedance state to an invalid data state prior to the specified access time.

Figure 5. Read-Modify-Write Cycle Timing

PRODUCT PREVIEW

PARAMETER MEASUREMENT INFORMATION



NOTE A: All \overline{xCAS} must be high.

Figure 9. \overline{RAS} -Only Refresh Timing

PRODUCT PREVIEW

PARAMETER MEASUREMENT INFORMATION

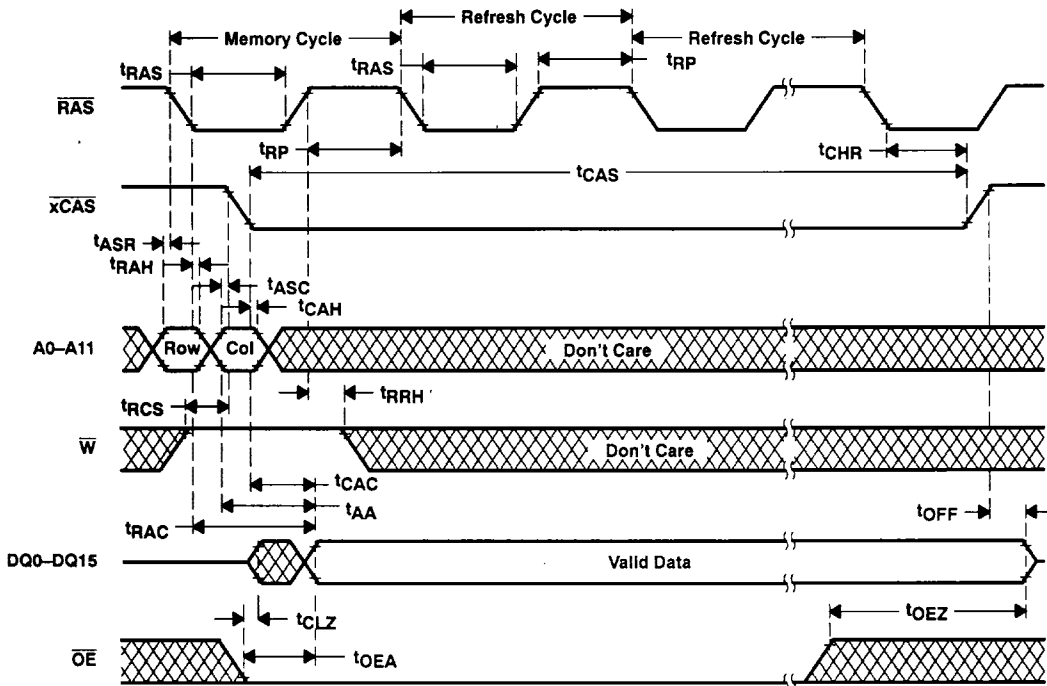
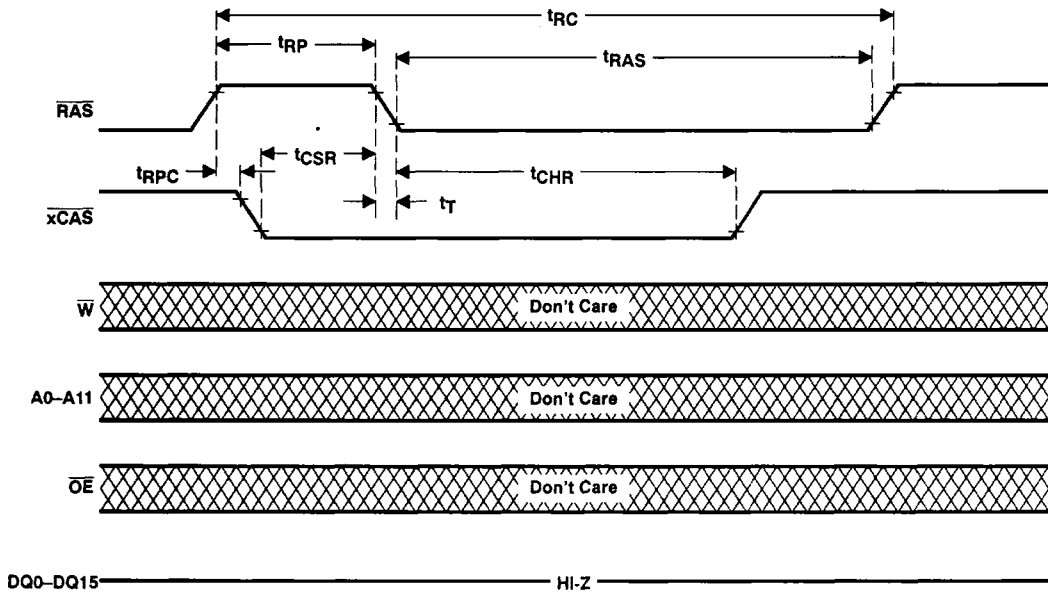


Figure 10. Hidden Refresh Cycle Timing

PRODUCT PREVIEW

PARAMETER MEASUREMENT INFORMATION



NOTES: A. Any \overline{xCAS} may be used.

Figure 11. Automatic (\overline{CAS} -Before- \overline{RAS}) Refresh Cycle Timing

PRODUCT PREVIEW

PARAMETER MEASUREMENT INFORMATION

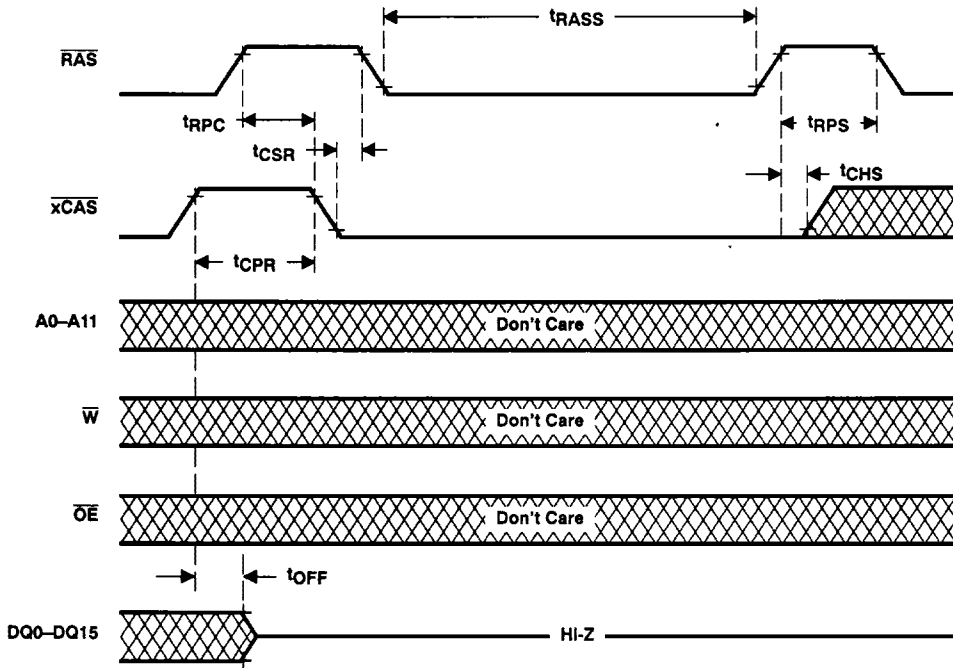


Figure 12. Self Refresh Timing

PRODUCT PREVIEW

device symbolization

